

# **Lecturer 4, 5**

# **Pulse Code Modulation**

Prepared by Prof

**Mahmoud Ahmed Attia Ali**

Department of Electronics and Communications

Faculty of Engineering

Tanta University

Egypt

For additional advice see  
Dale Carnegie Training®  
Presentation Guidelines



# Contents

- ❑ Generation of PCM Signal.
- ❑ Quantization and Coding.
- ❑ Companding Process.
- ❑ North American PCM T1 Standard.
- ❑ Intersymbol Interference ISI.
- ❑ Digital Signaling Formats.
- ❑ Spectral Efficiency.
- ❑ Differential Encoding.
- ❑ Synchronous and Asynchronous.

The image features a large, light blue diamond shape centered on a white background. The diamond is composed of two overlapping triangles. The left side of the image is decorated with a vertical bar consisting of a yellow rectangle at the bottom and a magenta rectangle at the top. The letters "PCM" are written in a bold, dark gray, sans-serif font across the center of the diamond. The letters have a slight 3D effect with a thin, lighter gray shadow offset to the right and bottom.

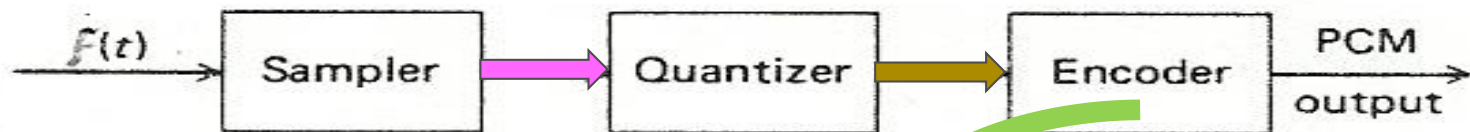
**PCM**

# Generation of PCM

Consists of three processes: ➡

- **Sampler:** Message signal  $f(t)$  is first sampled by a rate  $f_s > 2f_m$ .
- **Quantizer:** Sample values are then quantized to a certain levels.
- **Encoder:** Quantization levels are encoded into binary sequence.

# PCM Block Diagram



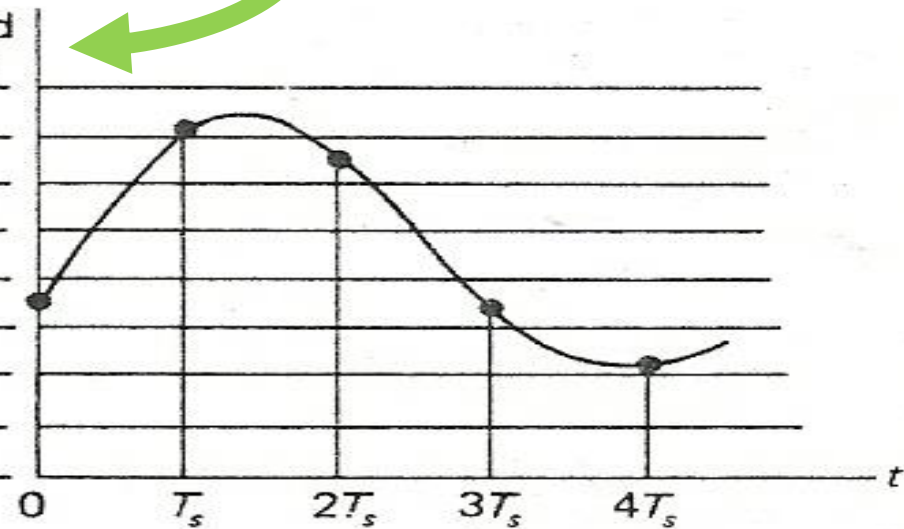
(a)

Quantization  
level  
number

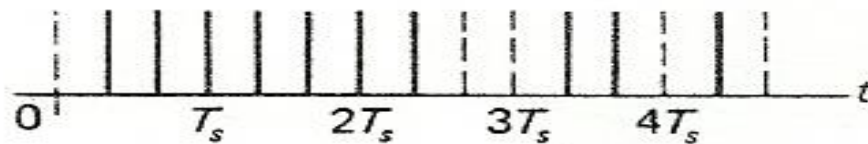
7
6
5
4
3
2
1
0

Encoded  
output

111
110
101
100
011
010
001
000

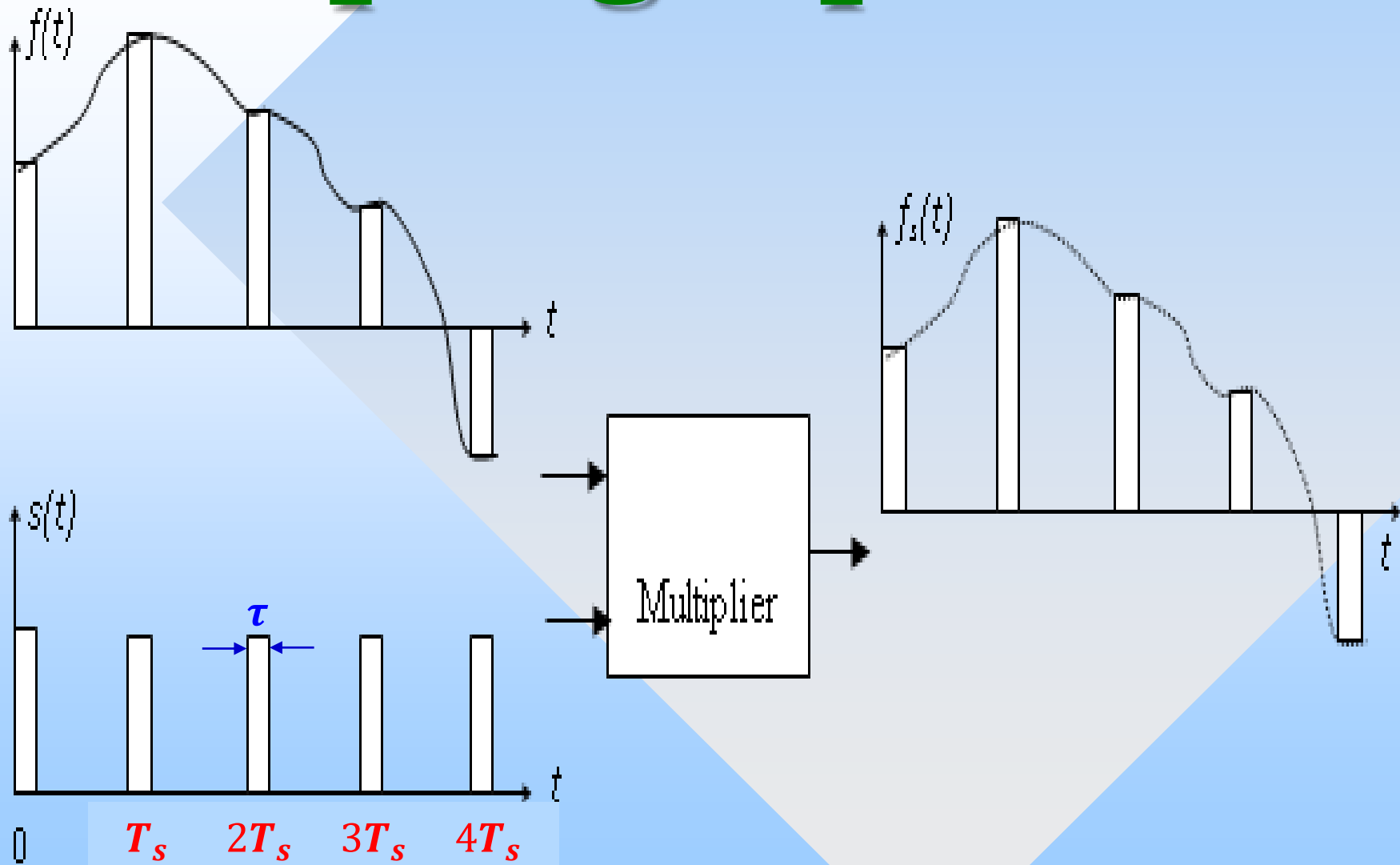


(b)

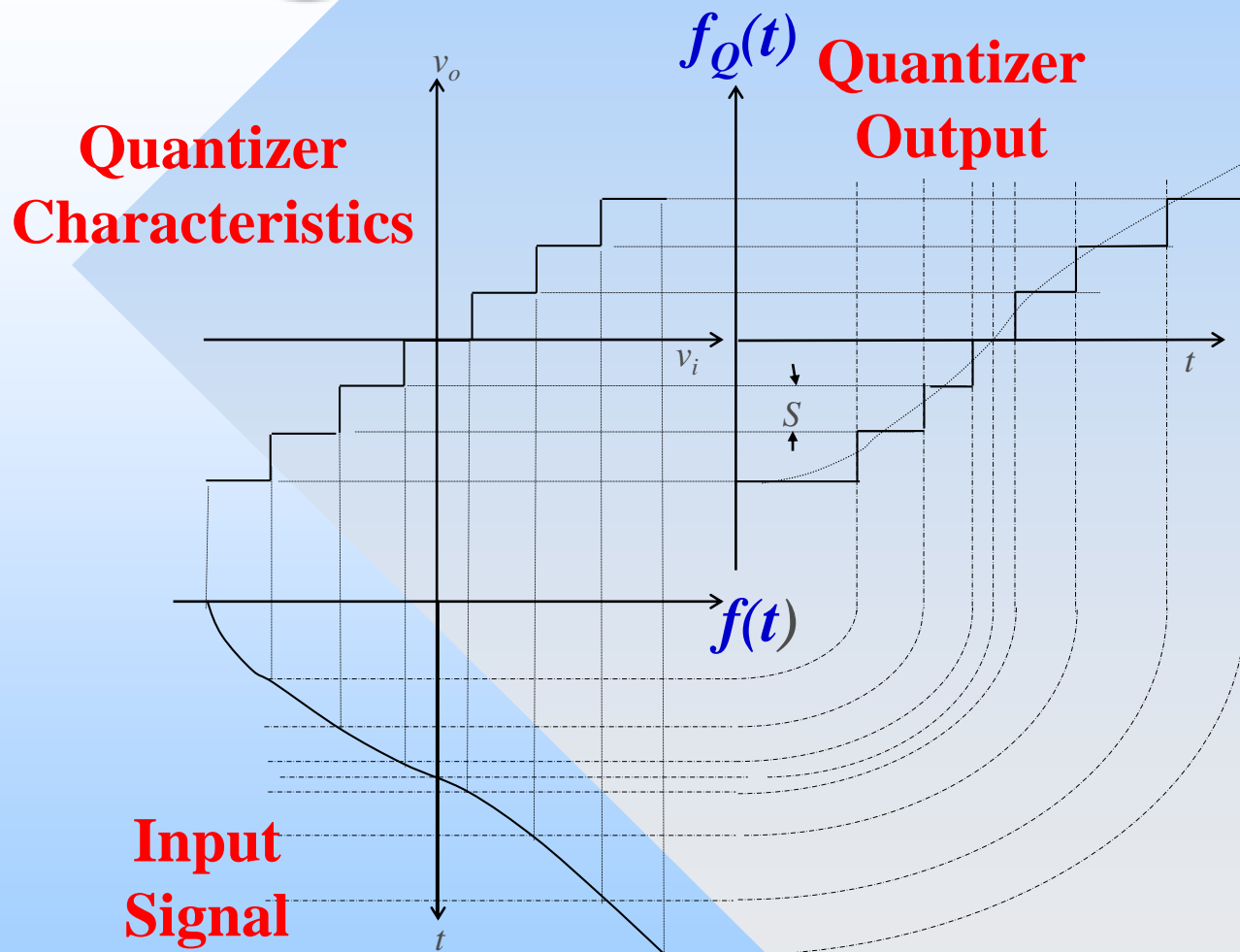


(c)

# Sampling Operation



# Quantization



# Quantization Error

- ❑ Difference between original signal  $f(t)$  and its quantized approximation  $f_Q(t)$
- ❑ Why termed as quantization noise?
  - ❑ Affects the signal amplitude.
  - ❑ May be added to or subtracted from the signal.
  - ❑ Expected average value is zero.
- ❑ Maximum value is  $\frac{1}{2}$  least significant

$$Q_{\text{maximum}} = \frac{v_{\text{LSB}}}{2}$$

- ❑ Upon reconstruction:
  - ❑ Added noise may be removed.
  - ❑ Sometimes errors take place.





$$\overline{e^2} = \frac{S^2}{12}$$

# Quantization Error

- ❑ Mean square quantization error voltage at quantizer output can be shown to be:

$$\overline{e^2} = S^2 / 12$$

- ❑ Quantization error voltage is:  $N_{ov} = \bar{e} = S / \sqrt{12}$

- ❑ Signal voltage at the quantizer output is:

$$S_{ov} = SM / 2$$

- ❑ Signal to quantization error voltage ratio:

$$\frac{S_{ov}}{N_{ov}} = \sqrt{3} M$$

- ❑ Signal to quantization error power ratio :

$$\frac{S_v}{N_o} = 3 M^2 = -4.8 + 20 \text{ Log } M \text{ dB}$$

# Quality of Quantization

- ❑ Quality of approximation improved by reducing the step size.
- ❑ Tests for speech indicate that:
  - ❑ 2 levels are understandable but quite noisy.
  - ❑ 8 or 16 levels are sufficient for a good intelligibility.
  - ❑ 128 or 256 are usually used to ensure high quality.
- ❑ Tests for color TV:
  - ❑ 64 levels gives only good color TV.
  - ❑ 512 levels is used for commercial color TV.

# Missed Signal Details

- ❑ In quantizing some details are lost.
- ❑ It is impossible to reconstruct the original.
- ❑ However, there is no need to transmit all signal details:
  - ❑ Ears and eyes in hearing and watching is limited:
    - ❑ Ear could not distinguish small distortion.
    - ❑ Eye has a limited resolution.
  - ❑ Also, due to noise, detector will not be able to distinguish fine variation.

# Probability of Level Error

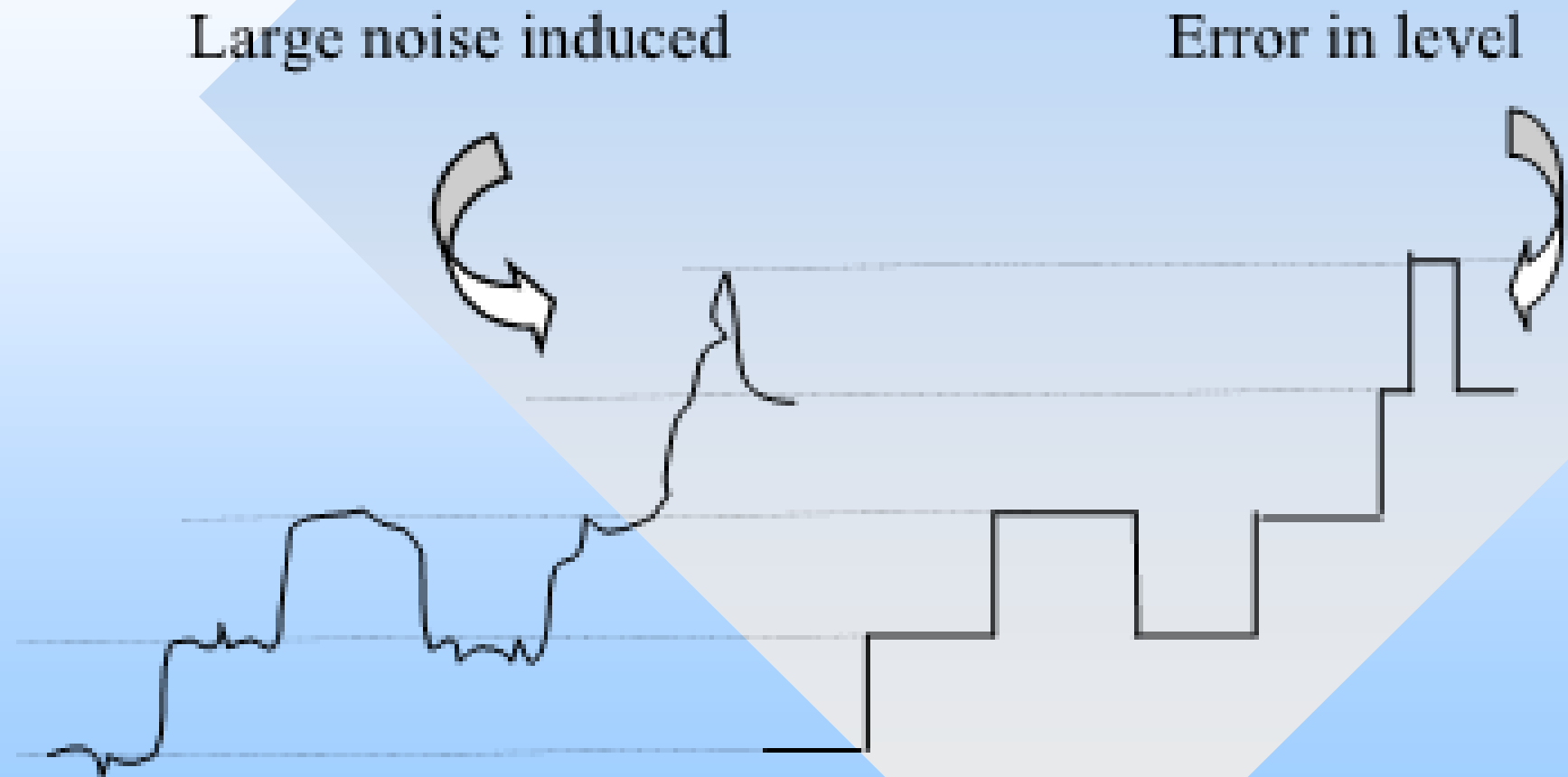


Fig.1.18: Noise Removing or Level Error

# Encoding

- ❑ A binary code where  $n$  equals 2.
- ❑ Number of quantization levels  $M$  is related to the number of bits per sample  $n$  as:

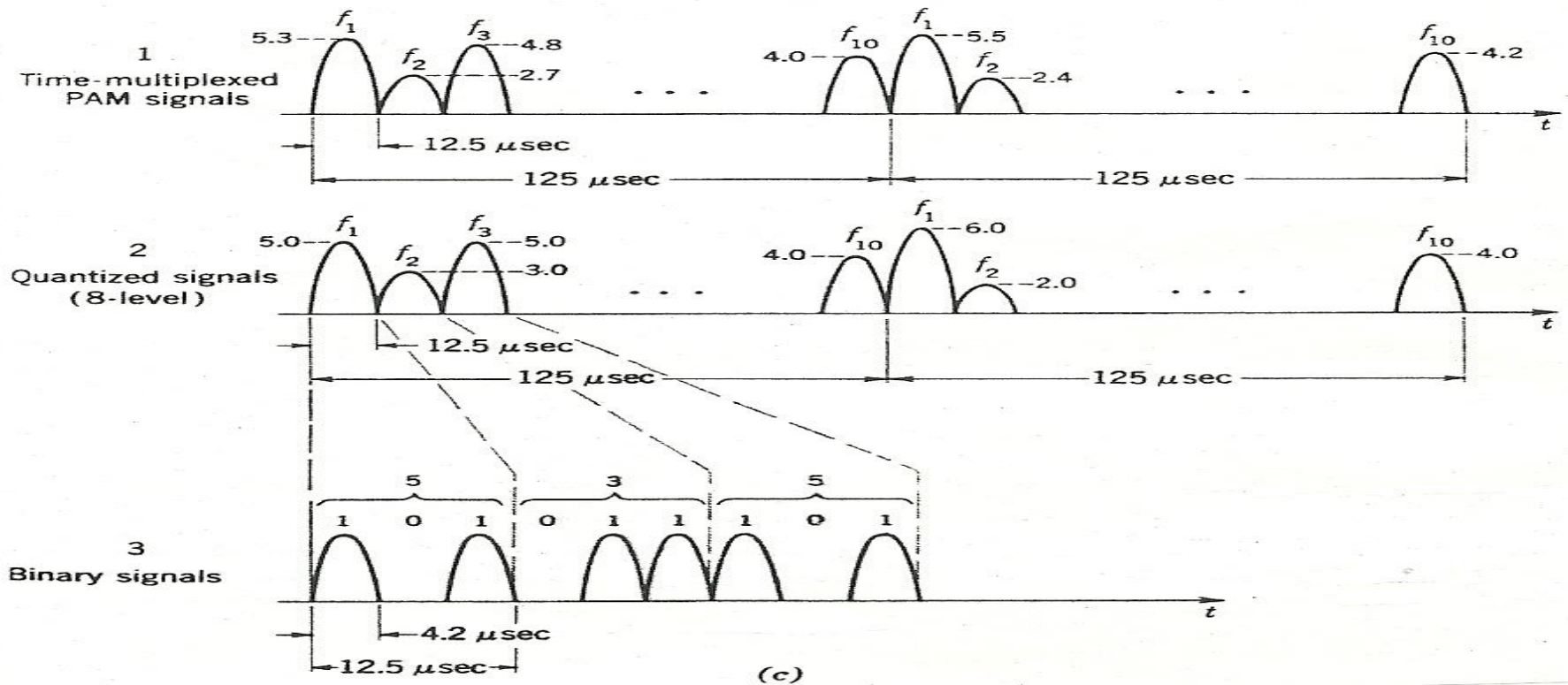
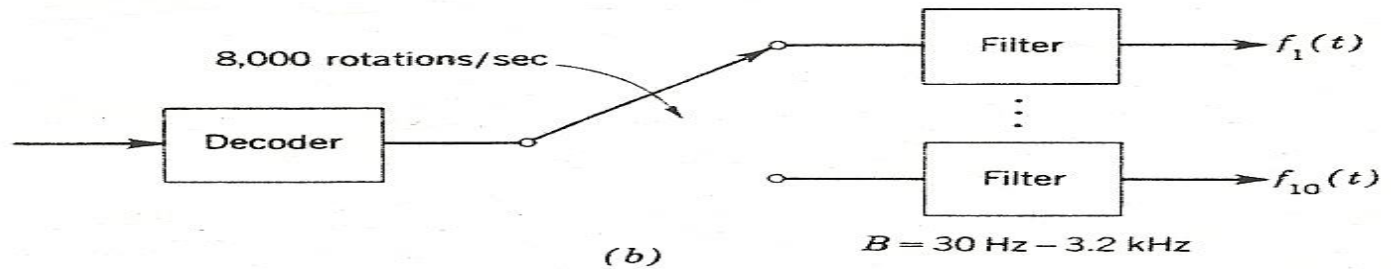
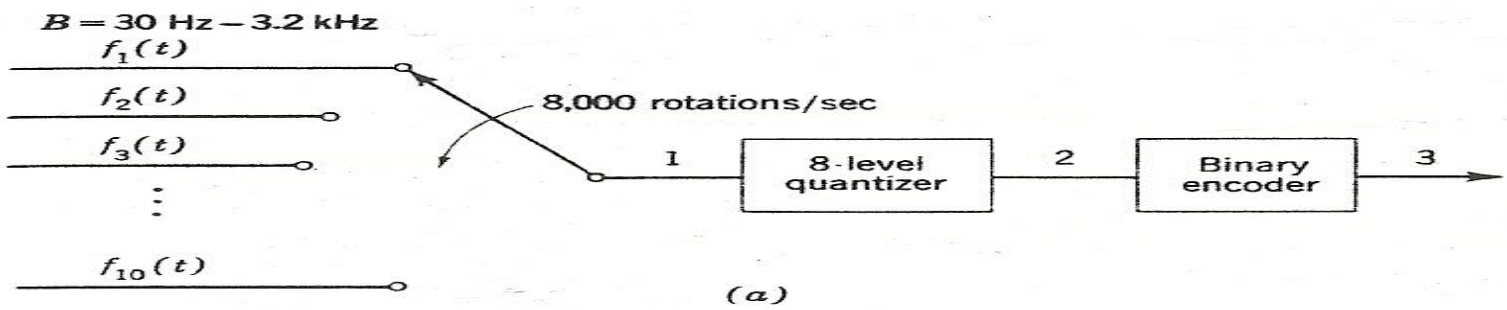
$$M = 2^n$$

- ❑ Generally, bandwidth for transmission of pulse train is inversely proportional to its width and depends on its shape.
- ❑ Therefore, bandwidth is roughly given by the reciprocal of the time slot.

# Example

- ❑ **Assume 10 channels PCM System**
  - ❑ Sampled PAM,
  - ❑ Quantized, and
  - ❑ PCM using 8 levels quantizer.
- ❑ **Required Bandwidths:**
  - ❑ PAM signal bandwidth is  $1/12.5 \mu\text{sec} = 80 \text{ kHz}$ .
  - ❑ Quantized PAM signal bandwidth is 80 kHz.
  - ❑ PCM signal bandwidth is  $1/4.2 \mu\text{sec} = 240 \text{ kHz}$ .





# Exercises

- ❑ Indicate the advantages of PCM systems when compared to PAM or Delta techniques.
- ❑ What is difference between the unipolar and the bipolar quantization?
- ❑ What is difference between mid-rise and mid-tread quantization processes?
- ❑ What is difference between the uniform and the nonuniform quantization?



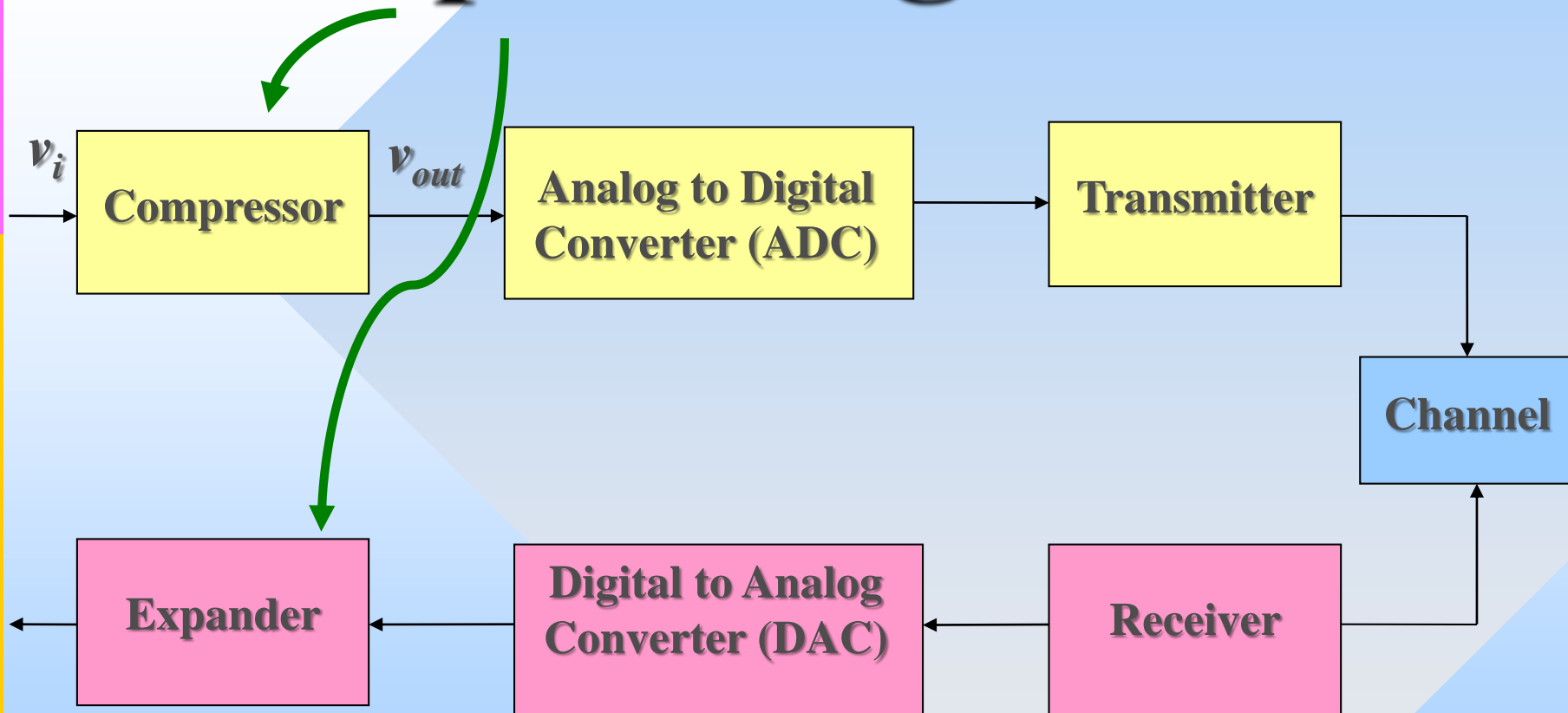


# Companding

# Reason of Companding

- ❑ Small signals will have a poorer signal to quantization noise ratio SQR than large.
- ❑ So, it is better to have moderate levels for both low and high variations of the signal:
  - ❑ Increase low signal levels to more moderate levels so that SQR could be increased.
  - ❑ Reduce high signal levels to more moderate levels in order to decrease high SQR.

# Componding Process



# Componding Laws

- ❑ Componding means compression and expansion.
- ❑ Compression by using special designed diodes prior to sampling circuit
- ❑ Whereas expansion is attained with diodes after the receiver LPF.
- ❑ Voice signal require a constant SQR over a wide dynamic range DR. This requires a logarithmic compression ratio. There are two methods:
  - ❑  $\mu$  Law Componding.
  - ❑ A Law Componding.

**$\mu$  Law**

**Companding**

**(USA and Japan)**

# $\mu$ Law Compression

$$v_o = \frac{\ln \left[ \left\{ 1 + \mu \left( \frac{v_i}{V_{i,max}} \right) \right\} \right]}{\ln[1 + \mu]} V_{o,max}$$

$V_{i,max}$ : Max amplitude of input signal before compression.

$V_{o,max}$ : Max amplitude of output signal after compression.

$v_i$ : Amplitude of input signal before compression for  $v_i > 0$

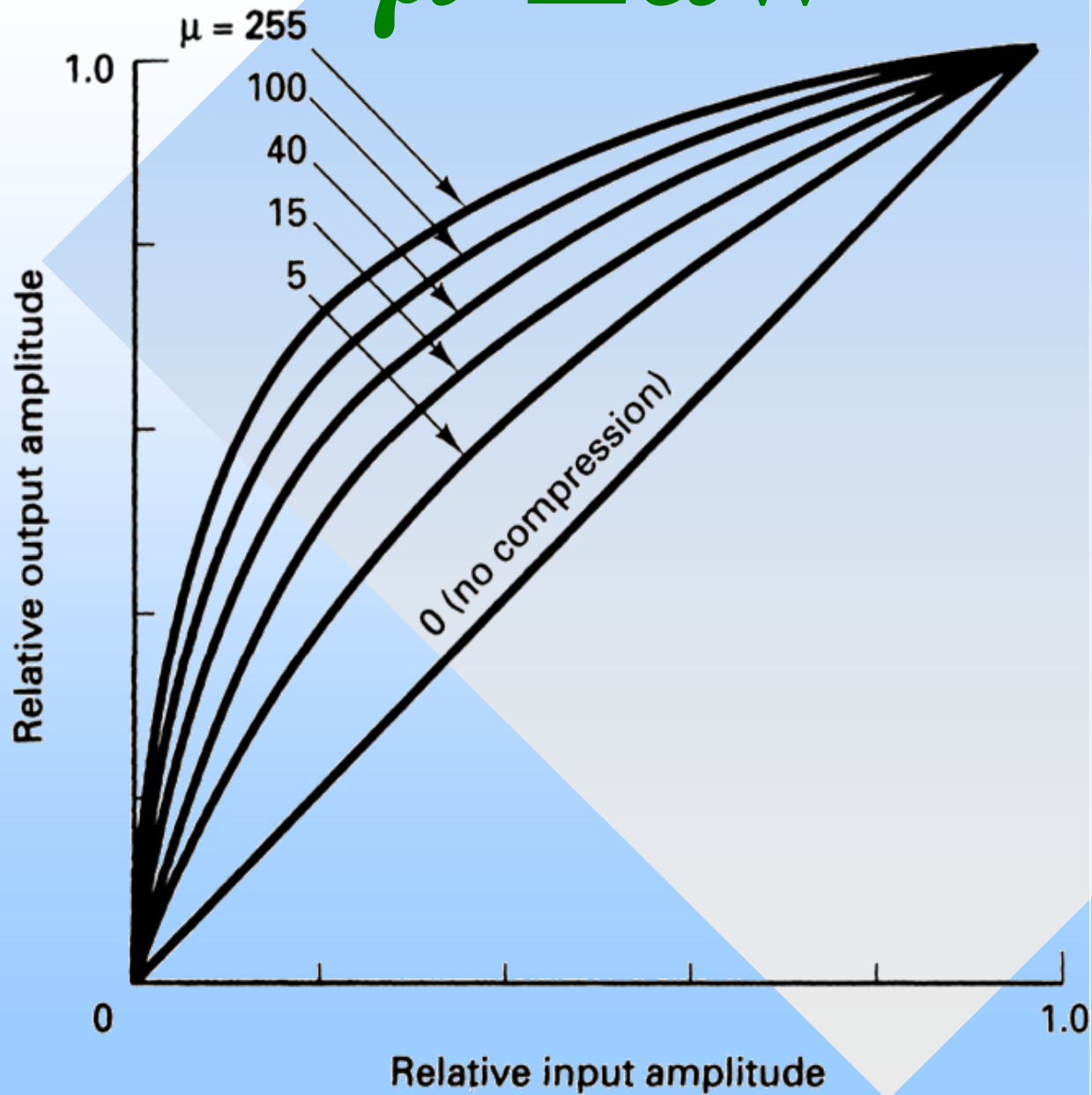
$v_o$ : Amplitude of the output signal after compression.

$\mu$ : is a measure for the amount of compression.

$\mu$ : determines the range of signal power over which SQR is relatively constant.

- For DR = 40 dB, 7 bit PCM code uses  $\mu = 100$ .
- For DR = 40 dB, 8 bit PCM code uses  $\mu = 255$ .

# $\mu$ Law





# $\mu$ Law Expansion

On reception, received signal will be expanded to satisfy linearity. Henceforth, received signal is:

$$v_{o,r} = \frac{V_{i,r,max}}{\mu} \left[ (1 + \mu) \left( \frac{v_{i,r}}{V_{o,r,max}} \right) - 1 \right] , v_{i,r} \geq 0$$

$V_{i,r,max}$ : Max amplitude of input received signal before expansion.

$V_{o,r,max}$ : Max amplitude of output received signal after expansion.

$v_{i,r}$ : Amplitude of input received signal before expansion.

$v_{o,r}$ : Amplitude of the output received signal after expansion.



# **A Law** **Companding** **(Europe CCITT)**

# A Law Compression

Compression characteristics is a true logarithm

$$v_o = \frac{\frac{Av_i}{V_{i,max}}}{1 + \ln A} V_{o,max}, \quad 0 \leq \frac{v_i}{V_{i,max}} \leq \frac{1}{A}$$

$$v_o = \frac{1 + \ln \left[ \frac{Av_i}{V_{i,max}} \right]}{1 + \ln A} V_{o,max}, \quad \frac{1}{A} \leq \frac{v_i}{V_{i,max}} \leq 1$$

$V_{i,max}$ : Max amplitude of input signal before compression.

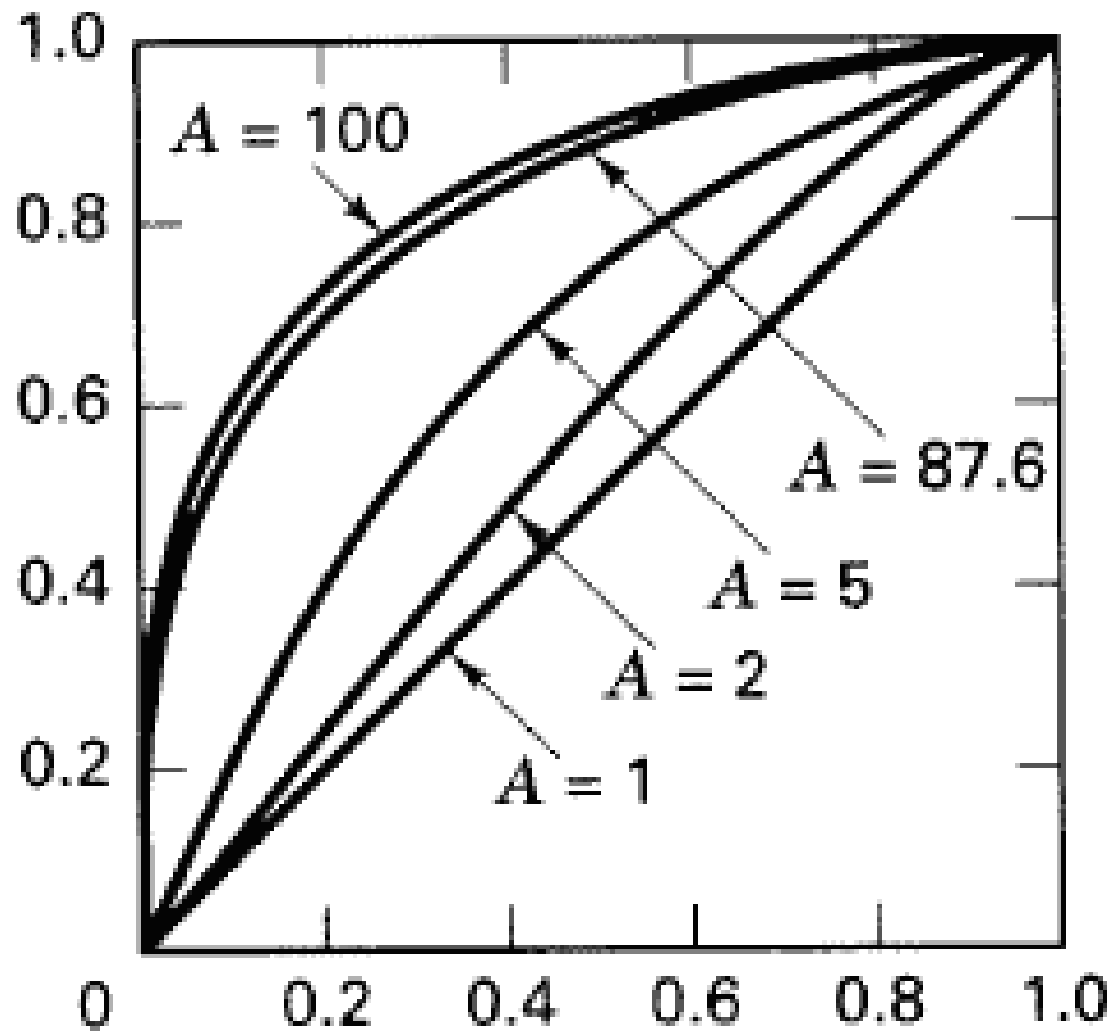
$V_{o,max}$ : Max amplitude of output signal after compression.

$v_i$ : Amplitude of input signal before compression.

$v_o$ : Amplitude of the output signal after compression.

Optimum value for voice transmission is  $A = 87.6$ .

# A Law



# Data Rate of PCM

- ❑ Speech signal for telephone has  $f_m = 4$  kHz.
- ❑ So, the sampling rate is  $2 \times 4 = 8$  kHz, that is 8000 samples/sec.
- ❑ Sampling interval  $1/8000 = 125 \mu$  sec/Frame
- ❑ Each sample is encoded into 8 bits/sample.
- ❑ So, the data rate for PCM signal is:  
$$R_{PCM} = (8000 \text{ samples/sec})(8 \text{ bits/sample})$$
$$= 64 \text{ k bits/sec}$$

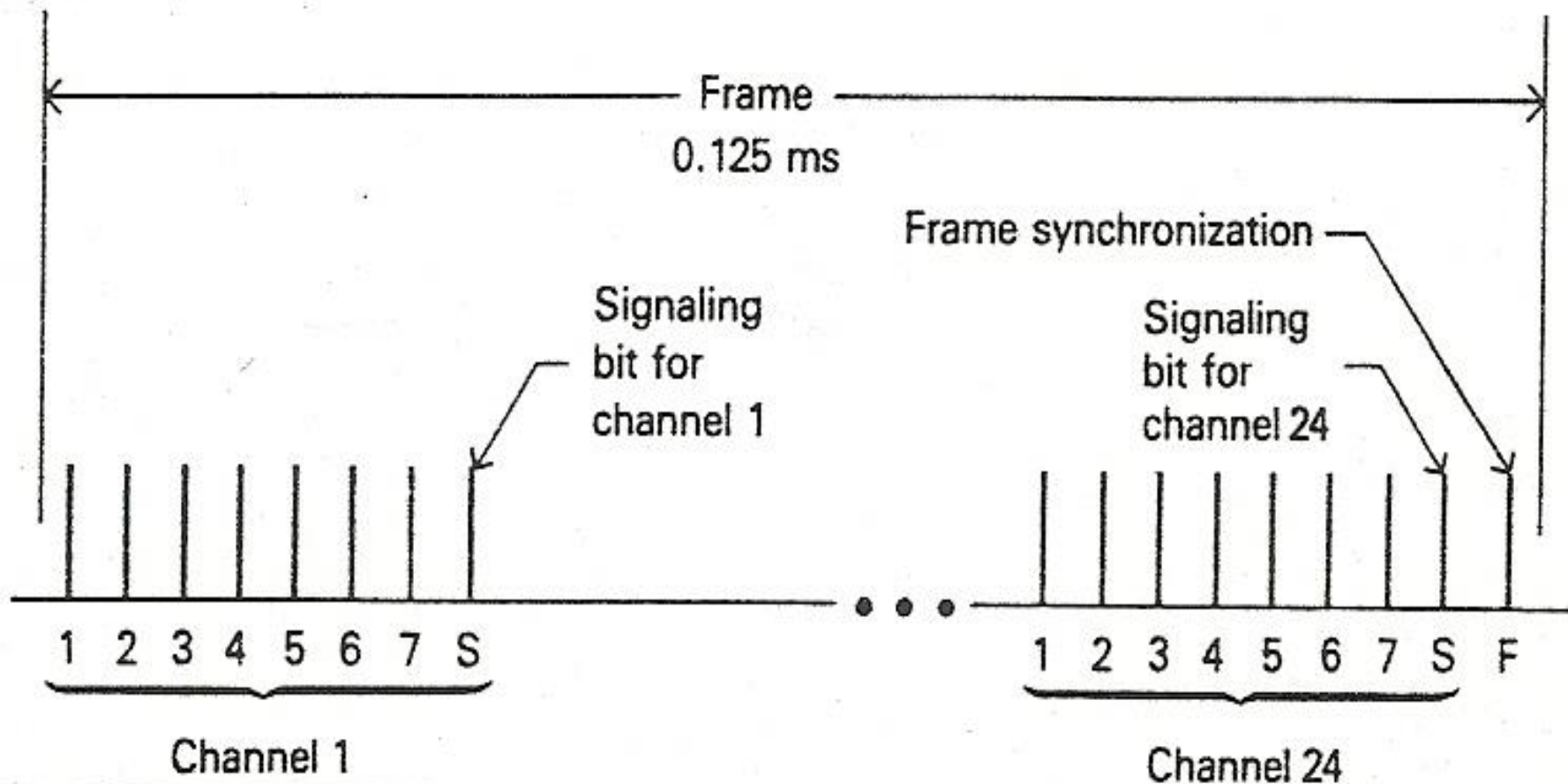
# North American PCM T1 Standard

- ❑ Bell System in USA introduced 24 channel PCM in 1960s for digital voice over short haul distances of 10 to 50 miles “T1”.
- ❑ T1 has found widespread adoption in US, Canada, and Japan.

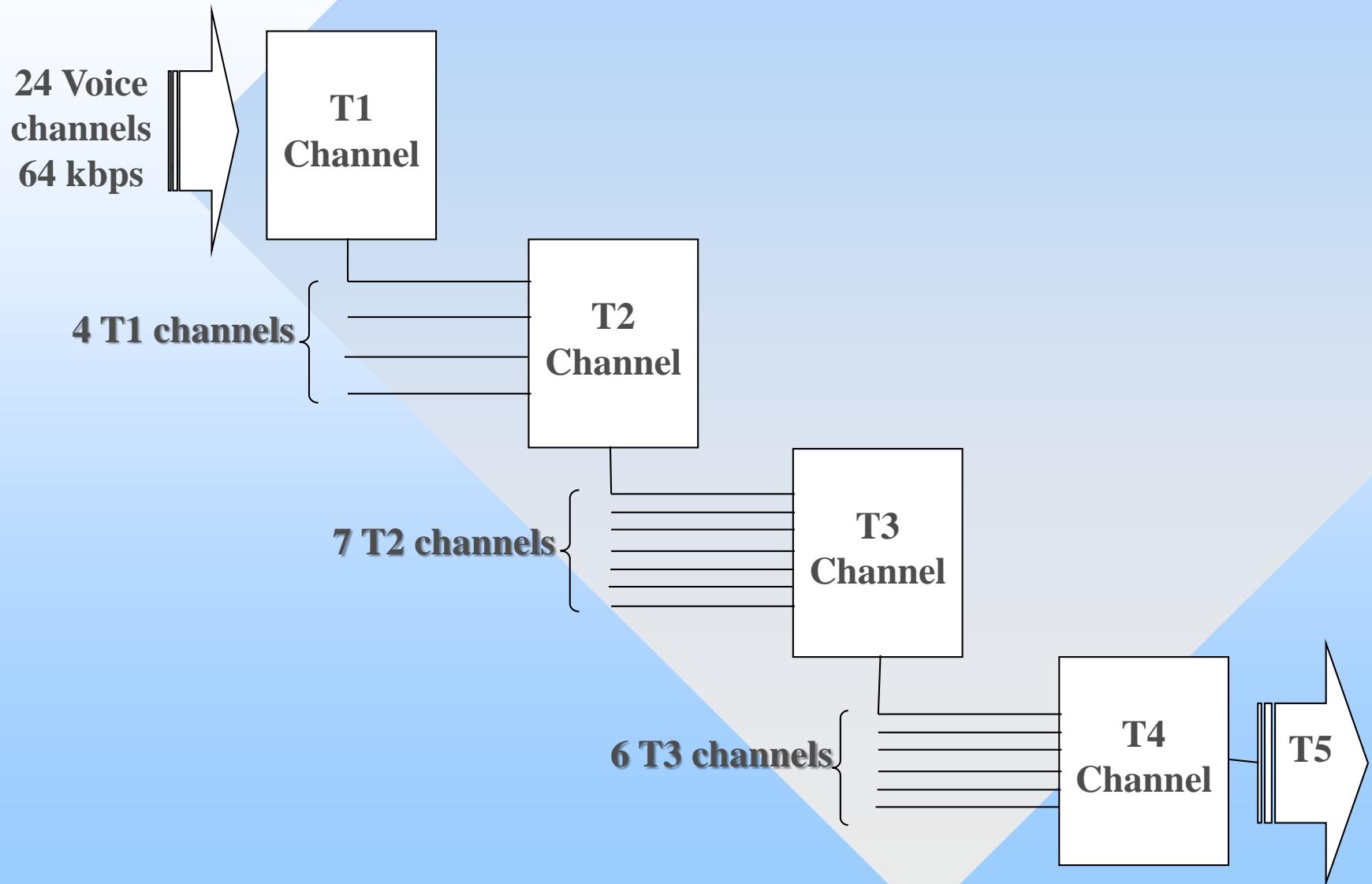
# PCM T1 Standard

- ❑ Early, it uses  $2^7 = 128$  quantization levels.
  - ❑ Each sample is quantized into 7 bits.
  - ❑ 1 bit for establishing calls (signaling).
- ❑ Recently,  $2^8 = 256$  levels have been adopted for quieter system with less distortion.
- ❑ 24 channels are time multiplexed, sampled, and coded into 8 bit PCM formats in addition to 1 bit for frame synchronization.
- ❑ The frame consists of  $24 \times 8 + 1 = 193$  bits.  
$$R_{T1} = (193 \text{ bits/Frame})(1\text{Frame}/125\mu\text{sec}) = 1.544 \text{ Mbps}$$

# North American PCM Standard for Short-Haul Telephone [T<sub>1</sub> System]



# Short to Long Haul PCM Systems





# Exercises

## Exercise 1.10

- ❑ Prove that the transmission data rate of T1 PCM system used in United States, Canada and Japan is given as 1.544 Mbps.

## Exercise 1.11

- ❑ Estimate the number of channels and the data rate for long haul T2, T3, and T4.

The image features a large, light blue diamond shape centered on a white background. The diamond is composed of two overlapping squares. The text "ISI" is written in a bold, dark gray, sans-serif font, centered within the diamond. The letters have a slight 3D effect with a drop shadow.

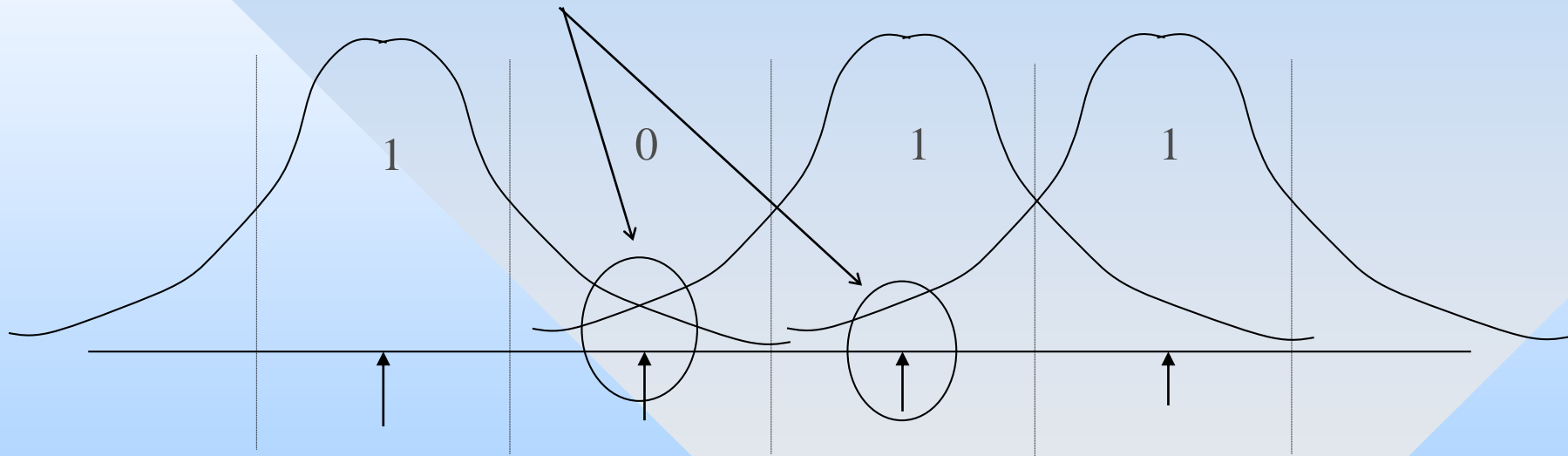
**ISI**

# Intersymbol Interference 'ISI'

- ❑ **First;** Bandwidth of PCM channel is restricted, waveforms will be distorted in a way analogous to crosstalk in PAM.
- ❑ **Second;** If digital pulses modulate a carrier for transmission over long distances, pulse shaping is a must.
- ❑ This causes pulses to spread out as they transverse channel and overlap into adjacent.

# ISI

## Inter-symbol Interference



# Crosstalk and ISI

- ❑ In un-quantized PAM, adjacent time slots are often associated with different message, and crosstalk is appropriate.
- ❑ In PCM adjacent bits are generally symbols in the code representation of a single quantized sample, hence the term inter-symbol interference.

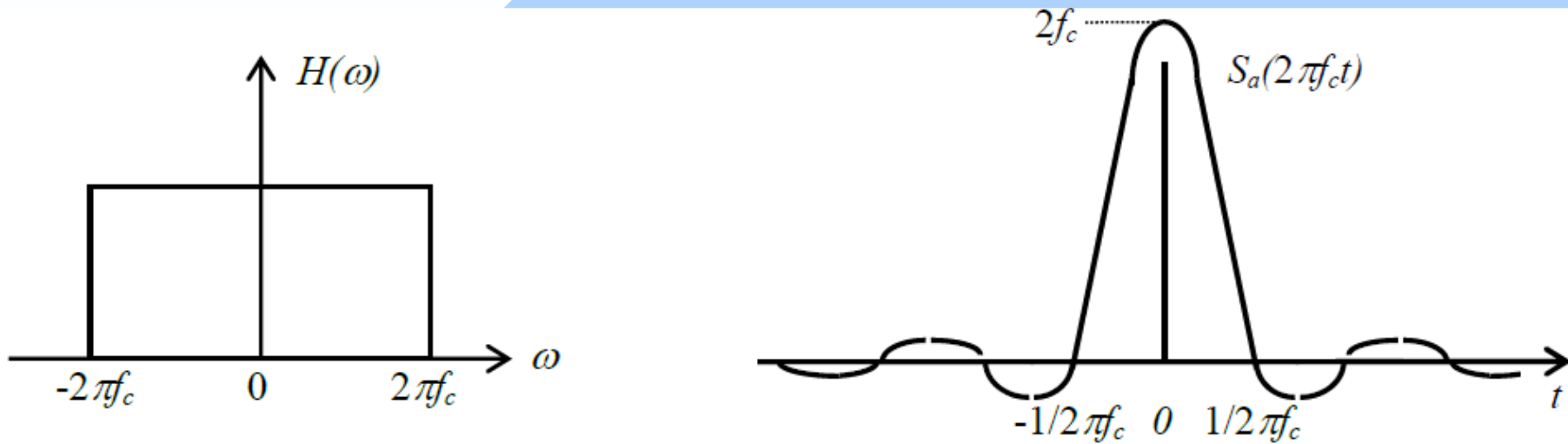
# Minimizing ISI

- ❑ Widening the transmission bandwidth as much as desired may minimize intersymbol interference.
- ❑ Instead, one could seek a way of purposely designing:
  - ❑ Signal wave-shapes and
  - ❑ Transmission filters used

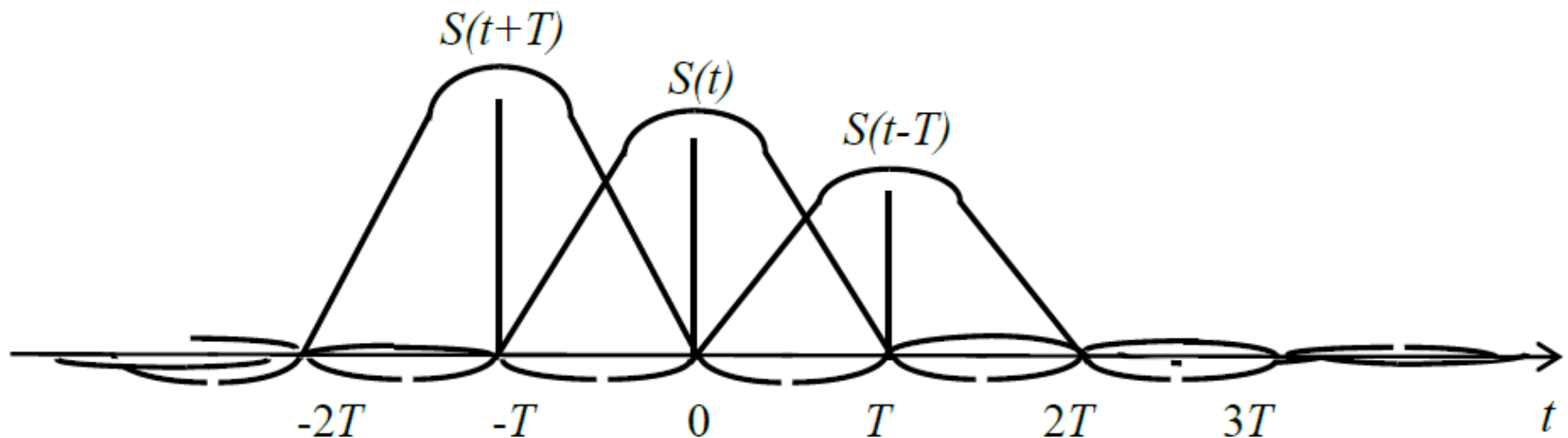
# Sampling Ch<sup>s</sup>

- ❑ One wave-shape producing zero ISI is the sampling function:  $(\sin 2\pi f_c t) / 2\pi f_c t$
- ❑ Is impulse response of ideal LPF at  $f_c$ :
- ❑ Pulse goes zero at multiples of  $1/2\pi f_c$
- ❑ If sample interval is chosen  $1/2\pi f_c$  adjacent pulses will not interfere

# Pulse Providing Zero ISI



## Sequence of Digital Pulses with Zero ISI





# Practical Difficulties

- ❑ It implies the overall ch<sup>s</sup> between the transmit and receive is ideal LPF.
- ❑ Require precise synchronization. If timing at receiver varies from exact synchronization, zero ISI condition disappears, the tails of all adjacent pulses may add up.
- ❑ Also some timing jitter will be present even with most sophisticated synchron.

# Raised Cosine

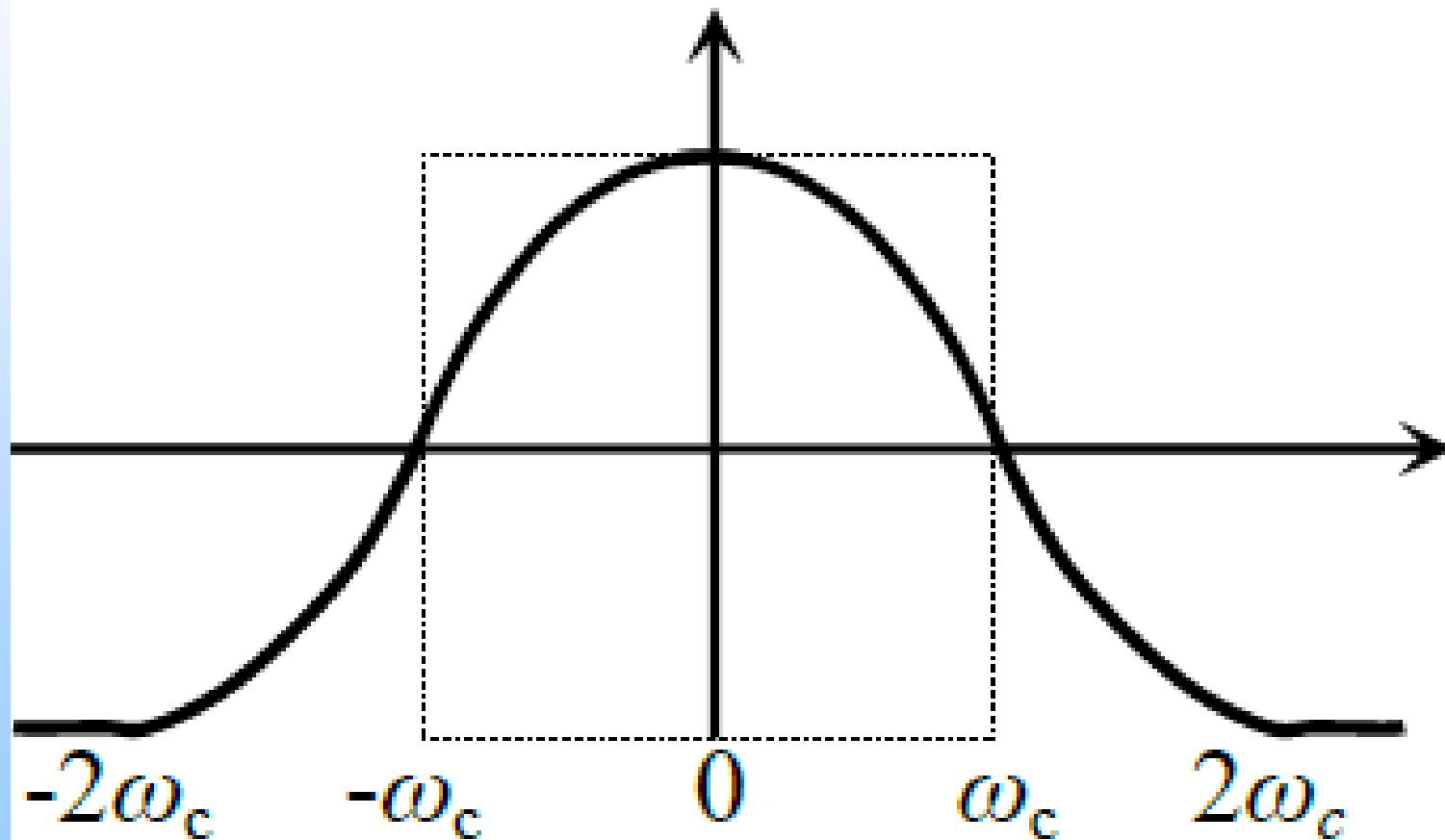
- ❑ If pulse has odd symmetry about LPF cutoff, impulse response retains property of zeros at uniformly spaced interval.
- ❑ An example is the raised cosine:

$$H(\omega) = \frac{1}{2} \left( 1 + \cos \frac{\pi \omega}{2 \omega_c} \right) \quad \& \quad \omega \leq 2 \omega_c$$
$$= 0 \quad \text{elsewhere}$$

- ❑ To show symmetry, let  $\omega = \omega_c + \Delta\omega$ :

$$H(\omega) = \frac{1}{2} \left( 1 + \cos \left\{ \frac{\pi}{2} + \frac{\pi}{2} \frac{\Delta\omega}{\omega_c} \right\} \right) = \frac{1}{2} \left( 1 - \sin \frac{\pi}{2} \frac{\Delta\omega}{\omega_c} \right)$$

# Cosine Odd Symmetry



# Impulse Response

- ❑ The impulse response of a filter with the characteristics shown is given by:

$$h(t) = \frac{\omega_c}{\pi} \frac{\sin \omega_c t}{\omega_c t} \left( \frac{\cos \omega_c t}{1 - (2\omega_c t / \pi)^2} \right)$$

- ❑ It has the  $(\sin x)/x$  term multiplied by an additional factor decreases with time.
  - ❑  $(\sin x)/x$  ensures zero crossing as for LPF.
  - ❑ Additional factor reduces the tails of pulses so that it becomes insensitive to timing jitter.



# **Digital Signaling Formats**

# Binary Line Coding

- ❑ Binary 1's and 0's may be represented in various signaling formats or line codes.
- ❑ The most popular are:
  - ❑ Unipolar Nonreturn to Zero, UNZ.
  - ❑ Unipolar return to Zero, URZ.
  - ❑ Bipolar Non return to Zero, PNZ.
  - ❑ Bipolar Return to Zero, PRZ.
  - ❑ Manchester

# Desirable Properties of Line Codes

- ❑ • Self-synchronization.
- ❑ • Error detection capability:
- ❑ • Low probability of error:
- ❑ • Suitable spectrum for channel:
- ❑ • Transparency:



# Self-synchronization

- ❑ There is enough timing information built into the code,
- ❑ So bit synchronizer can be designed to extract the clock signal.
- ❑ So, long series of binary 1's and 0's should not cause a problem in time recovery.

# **Error detection capability**

**The ability of addition of the channel encoders and decoders**

**Or incorporating them into the line code.**

# Low probability of error

- ❑ Receivers can be designed to recover data with low probability of error when the input is corrupted by
- ❑ Noise or
- ❑ Intersymbol interference ISI.

# Suitable Spectrum for Channel

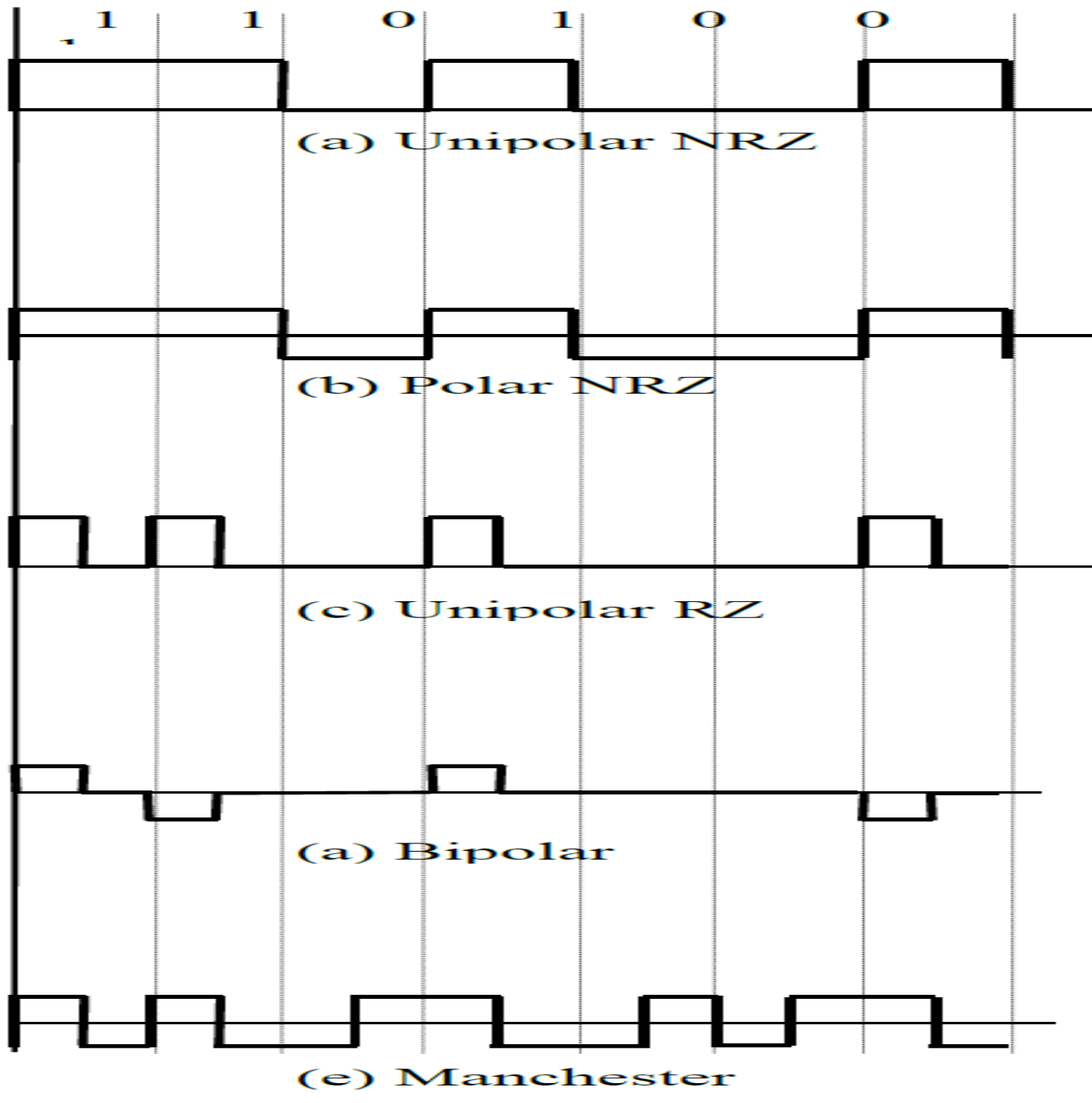
- ❑ With ac coupled, the power spectral density of line code should be negligible at frequencies near zero.
- ❑ Also, the signal bandwidth needs to be sufficiently small compared to channel bandwidth to minimize ISI.

# Transparency

- ❑ Every possible sequence of data is faithfully and transparently received.
- ❑ A code is not transparent if some sequence will result in a loss of synchronization at the receiver.
- ❑ For example, the bipolar format is not transparent since a string of zeros will result in a loss of clocking signal.

# Unipolar NRZ

- (on-off keying):
- The binary 1 is represented by a high level and
- a binary 0 by a zero level.
- High level does not return to zero during the binary 1 signaling interval.





# Polar NRZ

- **Simply NRZ:**
- **Binary 1's and 0's are represented by equal**
  - » **positive**
  - » **and negative levels.**

# Unipolar RZ

- The binary 1 is represented by a high level over half of the bit period and then returns to zero.

# Bipolar

- **Bipolar:**
- **The binary 1 is represented by alternatively positive or negative values over a half-bit period.**
- **The binary 0 is represented by a zero level.**

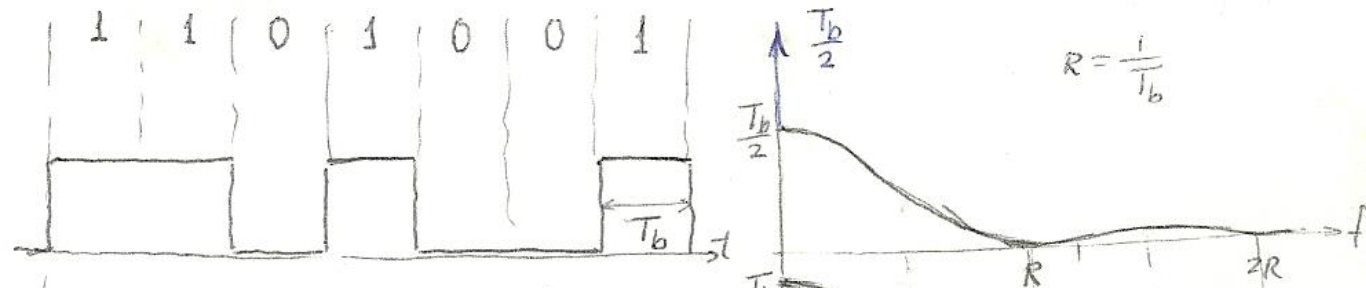
# Manchester

- ❑ Split phase encoding:
- ❑ Each binary 1 is represented by a positive half-bit period pulse followed by a negative half-bit period pulse.
- ❑ Similarly, a binary 0 is represented by a negative half-bit period followed by a positive half-bit period.

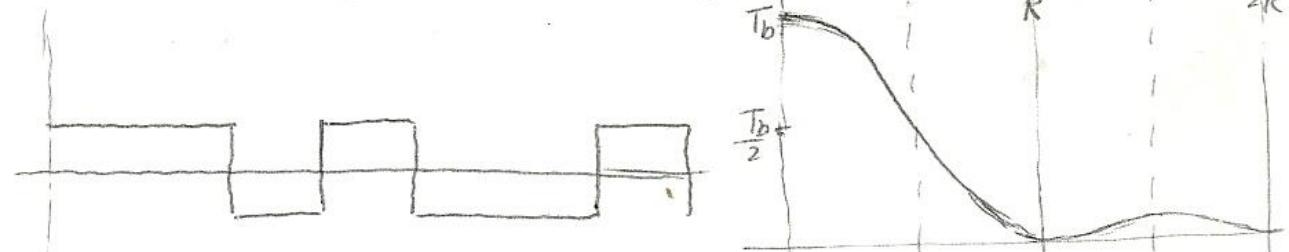
# **Power Spectra and Comparison**

# Power Spectra and Comparison

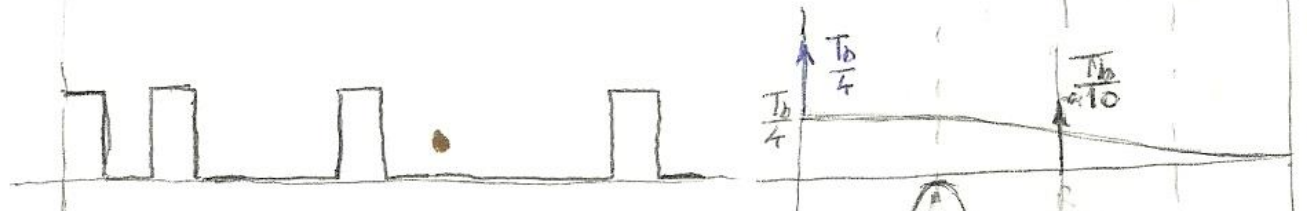
(a) Unipolar NRZ



(b) Polar NRZ



(c) Unipolar RZ



(d) Bipolar (AMI)



(e) Manchester



# Unipolar and Polar Signaling NRZ

- ❑ **Advantage:** They are easy to generate (TTL and CMOS circuits).
- ❑ **Disadvantage:** Waste of power due to the dc level so that dc coupled circuits are needed.
- ❑ Polar requires one supply although unipolar requires positive and negative supplies.
- ❑ Polar signaling has a superior bit error performance than other methods.



# Unipolar RZ

- ❑ In addition to the continuous power spectrum.
- ❑ It has discrete spectral lines at the odd multiples of the bit rate.
- ❑ It requires twice bandwidth of that of the NRZ codes.

# Bipolar

- ❑ Neither DC nor discrete spectra's.
- ❑ Bandwidth is that of NRZ.

# Manchester

- Neither DC nor discrete lines.
- Bandwidth is twice that of NRZ.

# Spectral Efficiency

- ❑ Number of bits per second of data that can be supported by each hertz of bandwidth.

$$\eta = \frac{R}{B} \quad (b/s/Hz)$$

- ❑ Unipolar, polar, and bipolar signals are twice as efficient as RZ or Manchester signals?
- ❑ Much greater efficiency can be achieved with multilevel signaling.
- ❑ Theoretical limit is determined by Shannon's?

$$\eta = \frac{C}{B} = \text{Log} \left[ 1 + \frac{S}{N} \right]$$

# Why Differential

- ❑ When serial data are passed through many circuits, waveform is often inverted (i.e., data complemented).
- ❑ This happen in twisted-pair channel just by switching the two leads at a connection point when polar signaling is used.
- ❑ For this reason differential coding is often employed

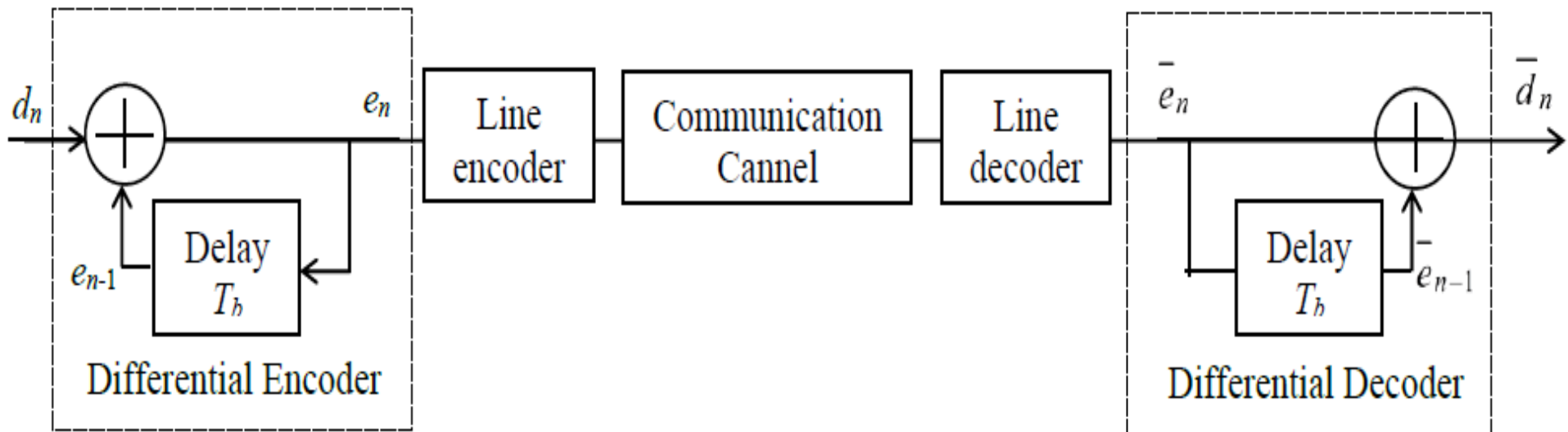
# Differential Encoding

- ❑ Each digit is obtained by comparing the present input bit with the past encoded bit.
- ❑ Binary 1 is encoded if the present bit and past encoded bit are of opposite state, and
- ❑ Binary 0 is encoded if the states are the same.
- ❑ Encoded differential data are generated by:

$$e_n = d_n \oplus e_{n-1}$$

$\oplus$  modulo 2 adder, or an exclusive OR gate

# Differential Coding System



# Differential Decoding

□ The received decoded data are given by:

$$\hat{d}_n = \hat{e}_n \oplus \hat{e}_{n-1}$$

$\oplus$  modulo 2 adder, or an exclusive OR gate



# **Synchronous** **and** **Asynchronous** **Lines**

# Synchronous System

- ❑ Each device is designed so that internal clock is relatively stable for a long period of time,
- ❑ And it is synchronized to a system master clock.
- ❑ Therefore, each bit of data is clocked in synchronism with the master clock.
- ❑ It requires a higher levels of synchronization to allow the receiver to determine the beginning and end of each block of data.
- ❑ Synchronizing signal may be:
  - ❑ Provided by a separate clocking line, or
  - ❑ Embedded in data signal (using Manchester)

# Example

US telephone industry plans to synchronize all of the digital lines with a master clock located in Hillsboro.

# Asynchronous System

- ❑ The timing is precise only for the bits within each character (or word).
- ❑ It is called start-stop signaling because each character consists of a start bit that starts the receiver clock and concludes with one or two stop bits that terminate the clocking.
- ❑ The receiver clock is started aperiodically.
- ❑ So, no synchronization with a master clock is required.

# Example

- ❑ Keyboard terminals are asynchronous sources.
- ❑ The complete character consists of 10 bits:
  - ❑ One start bit,
  - ❑ 7-bit of the ASCII code,
  - ❑ One parity bit, and
  - ❑ One stop bit (for  $R > 300$  bits/sec).
- ❑ In asynchronous TDM, different sources are multiplexed on a character-interleaved basis instead of interleaving on a bit-by-bit basis.

# Multiplexers Classified

- ❑ TDM to synchronous lines.
- ❑ TDM to quasi-synchronous lines:
  - ❑ Individual clocks of the input data sources are not exactly synchronized in frequency.
  - ❑ So there will be some variation in the bit rates from different sources.
  - ❑ Sometimes, the clock rates are not related by a rational number.
  - ❑ This requires stuff bits, which are dummy bits; 1's, 0's, or some alternating pattern, to compensate for the difference sources rates.
- ❑ TDM to asynchronous lines: It produces:
  - ❑ High-speed asynchronous output (no stuff bits required).
  - ❑ High-speed synchronous output (stuff bits is required).